



## **Preliminary FC64 57.8Gb/s PAM-4 FEC Capabilities 16-077v0**

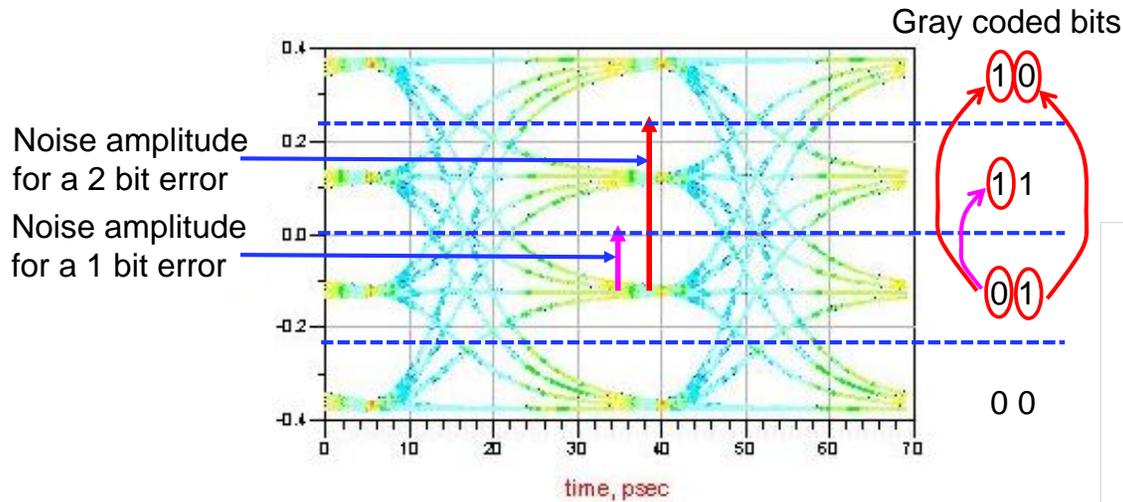
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February 3<sup>rd</sup>, 2016

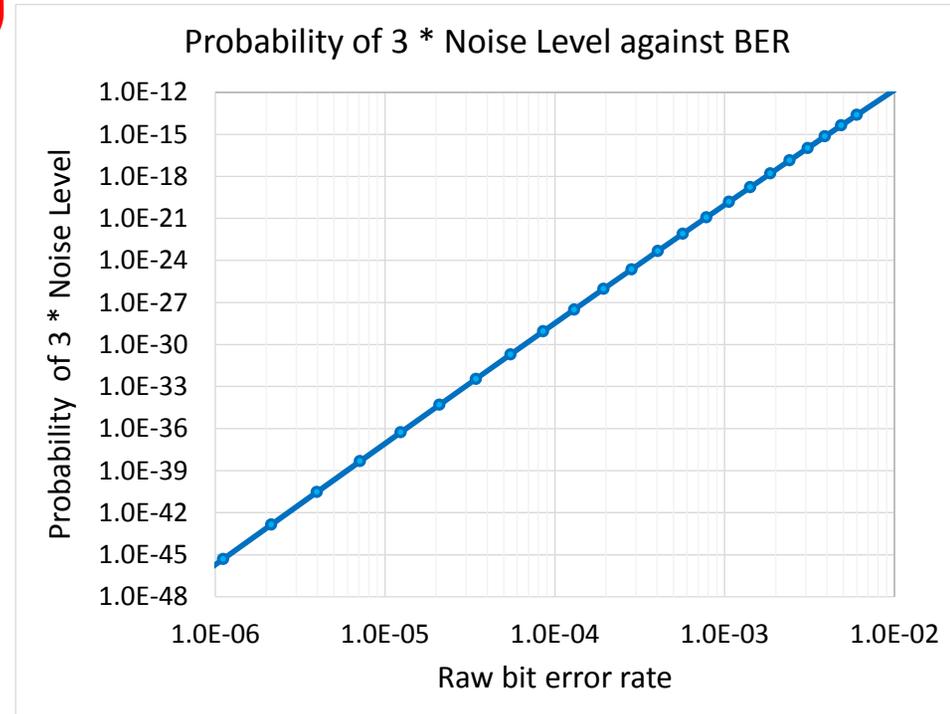
## Purpose of this work

- To determine if the (544,514,10) RS FEC allows any relaxation of the BER requirement of  $1e-6$  for the FC-56 standard optical link
- End to end link performance of  $1e-15$  BER is required after FEC correction
- It follows the flow in the IEEE presentation of Anslow\_3bs\_03\_0515

# Gray coded PAM-4 bit error rate probabilities



For a single bit in error the noise amplitude is  $\sim 0.5 * \text{level spacing}$   
 For a 2 bit error the noise amplitude is  $\sim 1.5 * \text{level spacing}$ ,  
 so =  $3 * \text{noise amplitude for a single bit error}$  assuming no  
 linearity error  
 Probability of 2 bit error is  $< 1e-20$  for  $1e-3$  BER and  $< 1e-45$   
 for  $1e-6$  BER, so 2 bit errors can be ignored for the rest of  
 this presentation



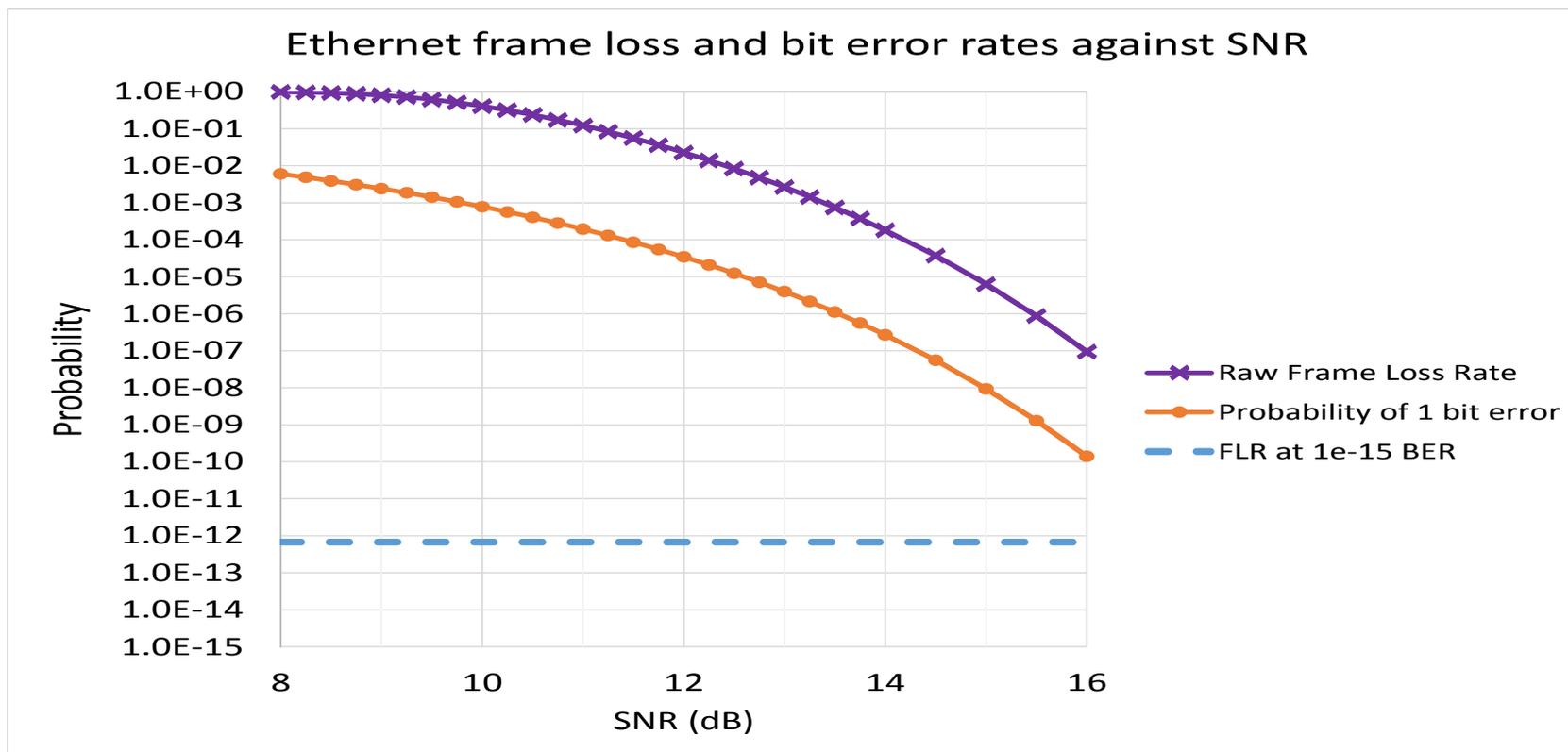
## Basic error probability theory

- The value of the Signal to Noise Ratio (SNR) used in the following curves is derived from the equation for a random noise induced error:

$$BER_{in} = \frac{1}{2} \operatorname{erfc} \left( \frac{SNR}{\sqrt{2}} \right)$$

- Where SNR is the signal to noise ratio expressed as a ratio and not in dB
- The minimum Ethernet frame length is  $(64 + 8 + 12) * 8 = 672$  bits and is used to correlate results with the Anslow presentation
- If the probability of a bit error is  $BER_{in}$  then the probability of a bit not being in error is  $(1 - BER_{in})$
- Therefore, if we assume no error propagation, then the probability of a frame of 672 bits having at least one error is  $1 - (1 - BER_{in})^{672}$  and is the Frame Loss Ratio (FLR)

# Random errors with no error propagation with no FEC

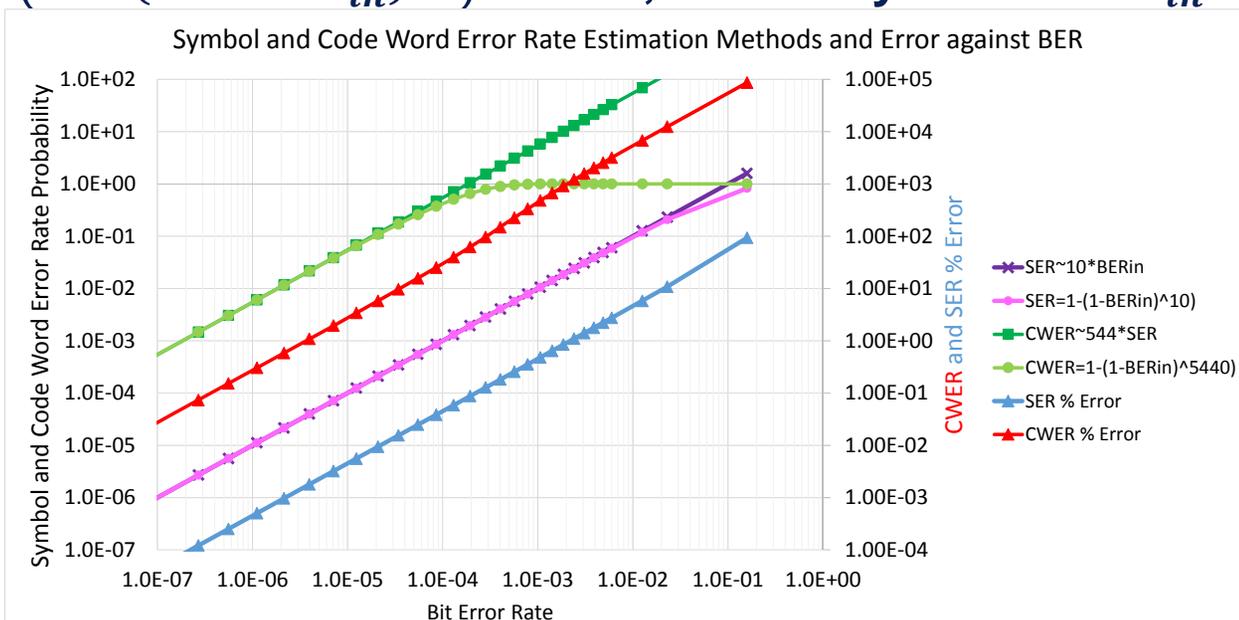


## (544,514,10) RS FEC error correction capability

- Block length  $n = 544$
- Message length  $k = 514$
- Symbol length = 10
- Hamming distance =  $n - k + 1 = 31$
- Erroneous symbols in a block that can be corrected =  $0.5 * (n - k) = 15$

# FEC mathematics with no error propagation

- For a BER of  $BER_{in}$  the probability of 1 bit not having an error is  $(1 - BER_{in})$
- The probability of a 10 bit symbol not having an error is  $(1 - BER_{in})^{10}$
- The probability of a 10 bit symbol having any number of errors is:  
 $(1 - (1 - BER_{in})^{10}) = SER$ , for a very small  $BER_{in}$  this reduces to  $10 * BER_{in}$



How small can we use the approximation?

As we have 544 symbols per code word the probability of no errors in a code word is

$$(1 - \text{Symbol Error Rate})^{544}$$

The simple approximation would be  $5440 * BER$  and only gives less than 1% error for  $BER < \sim 3e-6$  and  $> 100\%$  error for  $3e-4$  and above, so not good for the BER range we need to consider

SER = Symbol error rate

CWER = Code word error rate

## FEC mathematics with no error propagation – continued

- The probability of an error free symbol  $P(0) = 1 - SER$
- The probability of 544 error free symbols  $P(544, no\ errors) = (1 - SER)^{544}$
- The probability of one errored symbol  $P(1) = 1 - (1 - BER_{in})^{10} = SER$
- Number of ways of picking one errored symbol from 544 symbols = 544
- The probability of one errored symbol in 544 symbols  $P(1, 544) = 544 * SER$
- The probability of 543 symbols with no errors  $P(543, no\ errors) = (1 - SER)^{543}$
- The probability of one errored symbol and 543 non-errored symbols  

$$P(1\ error\ only) = 544 * SER * (1 - SER)^{543}$$
- The probability of two errored symbols in 2 symbols  $P(2) = SER^2$
- Number of ways of picking 2 symbols from 544 =  $\frac{544!}{(544-2)! * 2!}$
- Probability of 2 errored symbols in 544 =  $\frac{544! * SER^2}{(544-2)! * 2!}$
- The probability of 542 symbols with no errors  $P(542, no\ errors) = (1 - SER)^{542}$
- The probability of two errored symbols and 542 non-errored symbols  

$$P(2\ errors\ only) = \frac{544! * SER^2}{(544 - 2)! * 2!} * (1 - SER)^{542}$$

## FEC mathematics with no error propagation – continued 2

- The probability of  $n$  errored symbols  $P(n) = SER^n$
- Number of ways of picking  $n$  errored symbols from 544 =  $\frac{544!}{(544-n)!n!}$
- The probability of  $544-n$  symbols with no errors  
$$P(544 - n, \text{no errors}) = (1 - SER)^{544-n}$$
- Probability of  $n$  errored symbols in 544 =  $\frac{544! * SER^n}{(544-n)! * n!}$
- The probability of  $n$  errored symbols and  $544-n$  non-errored symbols  
$$P(n \text{ errors only}) = \frac{544! * SER^n}{(544-n)! * n!} * (1 - SER)^{544-n}$$

## FEC mathematics with no error propagation –continued 3

- The probability of up to 15 errored symbols

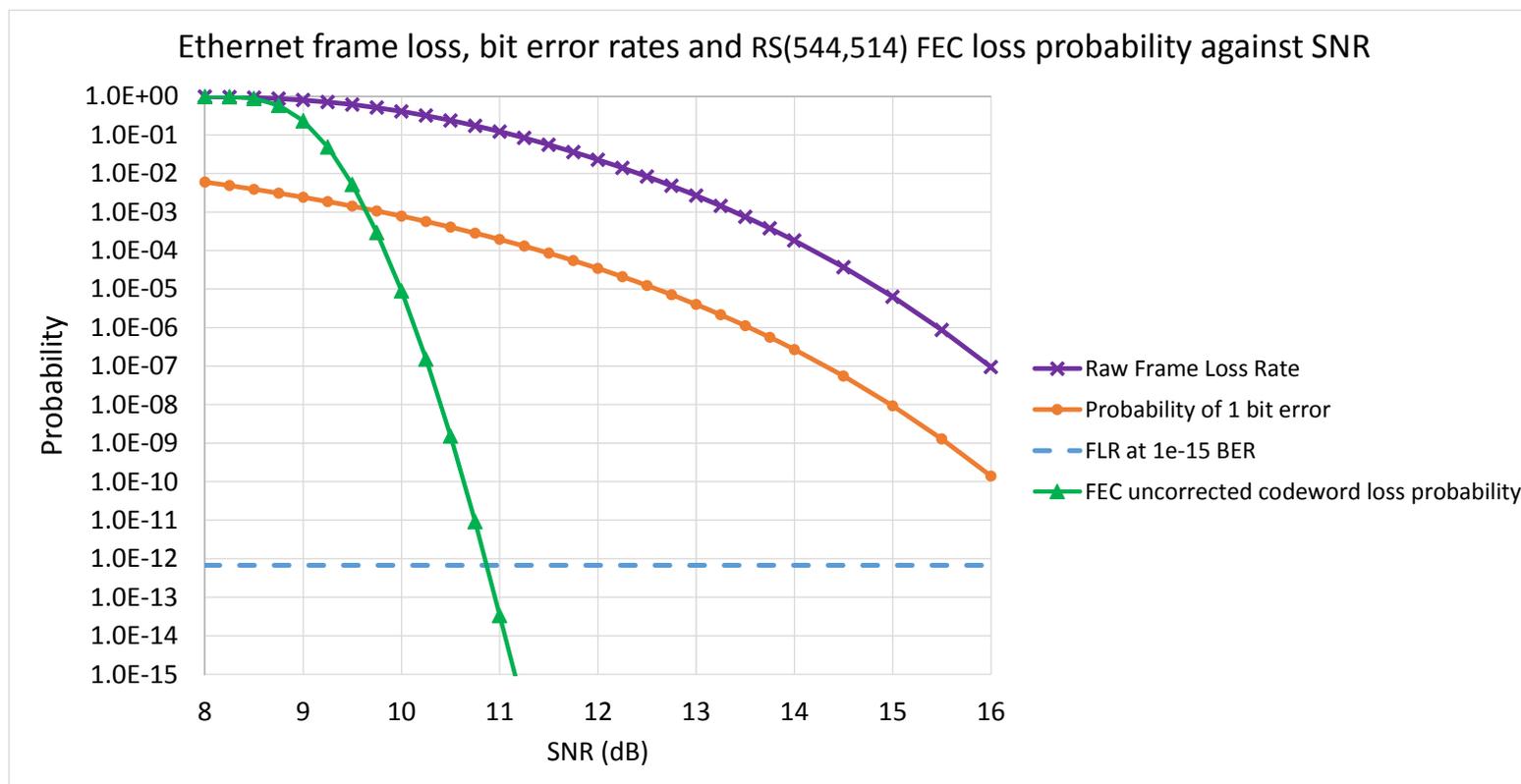
$$P(0\text{to}15\text{ errors}) = \sum_{n=0}^{n=15} P(n\text{ error only}) = \sum_{n=0}^{n=15} \frac{544! * SER^n}{(544 - n)! * n!} * (1 - SER)^{544-n}$$

- All of the 0 to 15 errored symbols cases will be corrected by the FEC, so the probability of an uncorrected error will be

$$P(FEC\text{uncorrected}) = (1 - P(0\text{to}15\text{ errors}))$$

$$P(FEC\text{uncorrectable}) = 1 - \sum_{n=0}^{n=15} \frac{544! * SER^n}{(544 - n)! * n!} * (1 - SER)^{544-n}$$

# Random errors with no error propagation with and without FEC

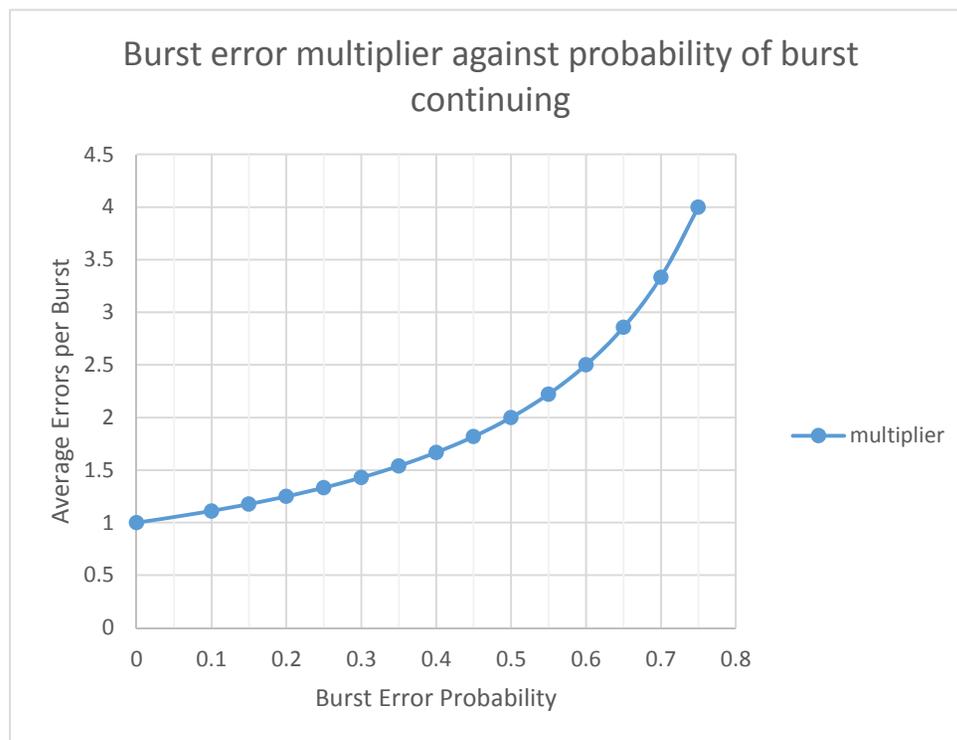


## Effect of Burst Error Probability on BER

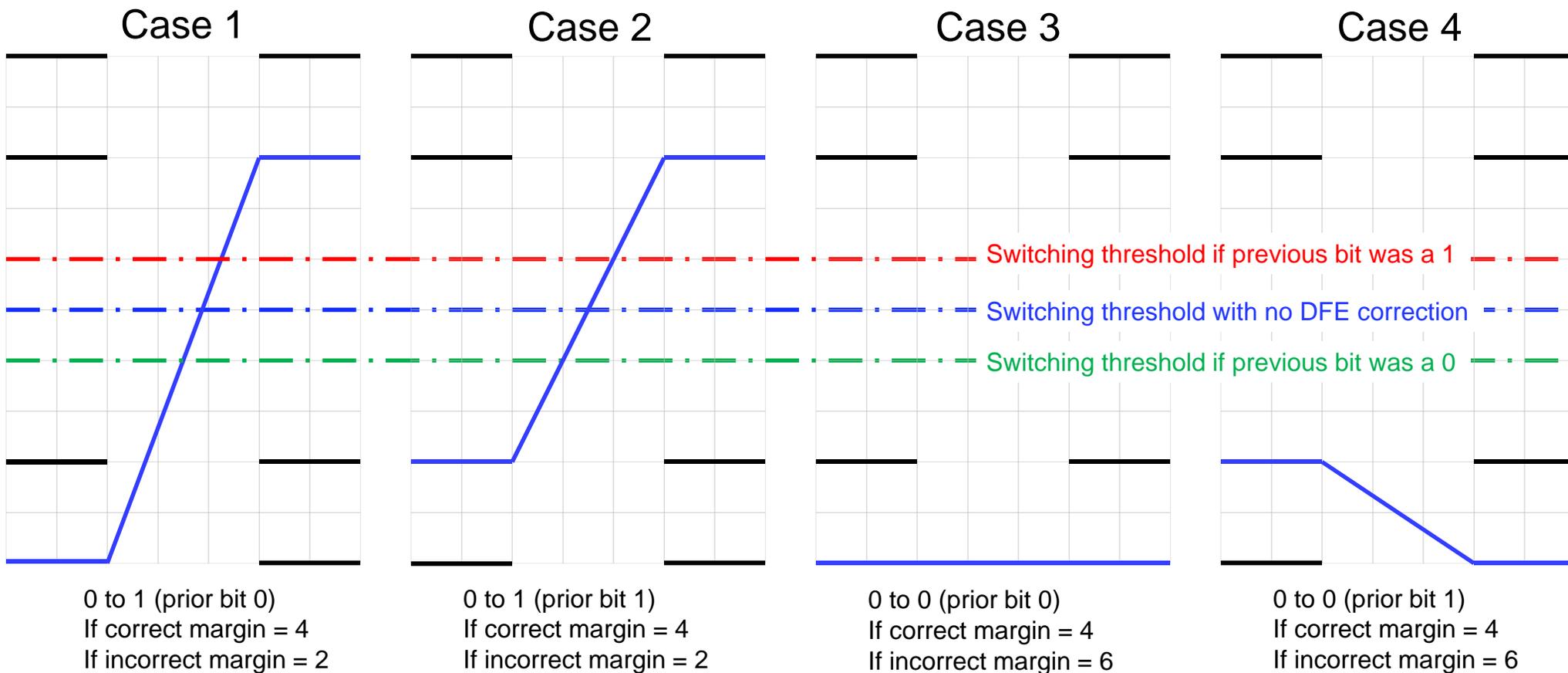
- The average number of errors in a burst is simply related to the probability of an error propagating (A)

$$\text{Average errors per burst} = \frac{1}{1 - A}$$

- However, with PAM-4 there are only 5 PAM-4 symbols per 10 bit FEC symbol
- For an error in the last bit of one FEC symbol with PAM-4 to propagate into the next symbol is very high, but to propagate to the following FEC symbol after that is also high and needs to be considered



# NRZ Error propagation options for 1 DFE tap only



In 50% of all cases the margin is reduced by  $2 * \text{DFE tap 1 value}$  and could cause an error to propagate

## Worst Case PAM-4 Error Propagation

Correct level	Perceived RX level		Errored bits	
	One higher	One lower	One higher	One lower
3 (1,0)	3 (1,0)	2 (1,1)	✓, ✓	✓, ✗
2 (1,1)	3 (1,0)	1 (0,1)	✓, ✗	✗, ✓
1 (0,1)	2 (1,1)	0 (0,0)	✗, ✓	✓, ✗
0 (0,0)	1 (0,1)	0 (0,0)	✓, ✗	✓, ✓

Two of these cases result in ✓, ✓ so the error will not propagate (25%)

This means a 75% chance of a PAM-4 error propagating

Therefore, average burst length = 4 PAM-4 symbols

Two of these cases result in errors in the first bit (25%), four cases in the second bit (50%)

## FEC Symbol Error Propagation for PAM-4 Errors

The probability of an error in the first PAM-4 symbol affecting the second PAM-4 symbol is 0.75, the third PAM-4 symbol is  $0.75^2$ , the fourth is  $0.75^3$ , the fifth  $0.75^4$  and the first PAM-4 symbol in the next FEC symbol is  $0.75^5$ .

This can be done for the second, third, fourth and fifth PAM-4 symbols being in error and gives the total probability of an error in one FEC symbol causing the next FEC symbol to also be in error of:

$$P_{err}(\text{next FEC symbol}) = \frac{(0.75^5 + 0.75^4 + 0.75^3 + 0.75^2 + 0.75^1)}{5} = 0.4576$$

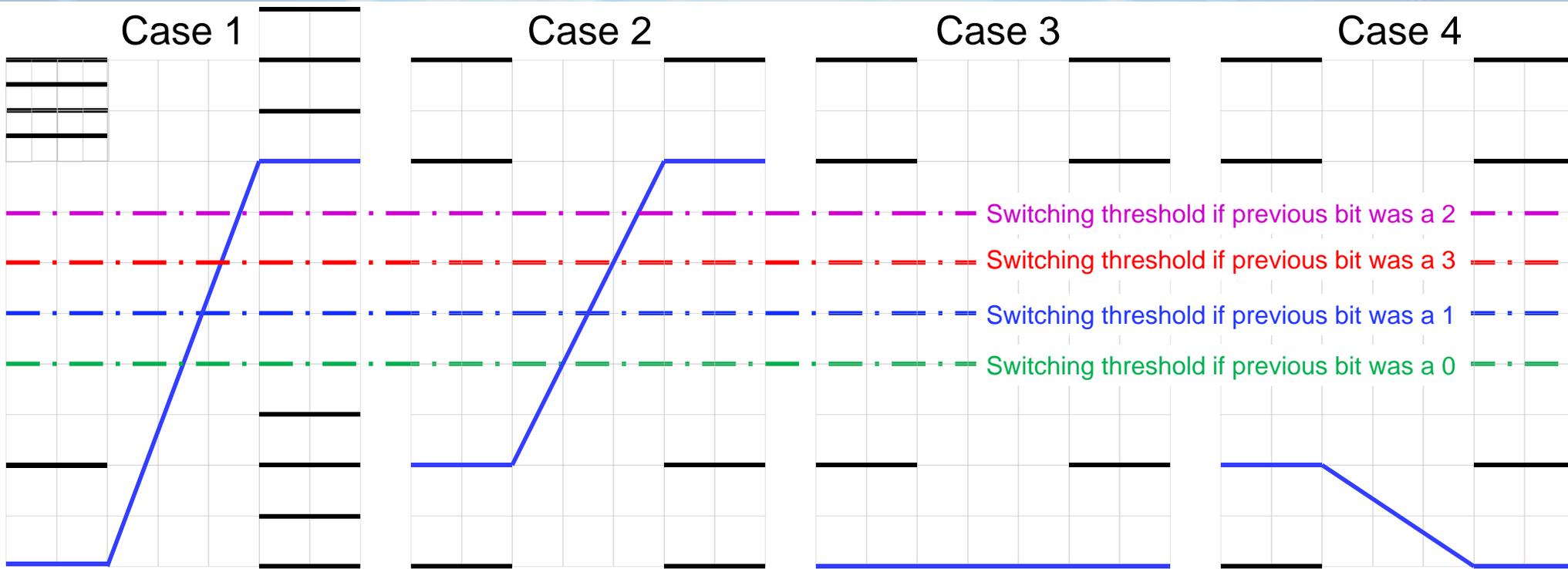
This gives a 46% chance of the next FEC symbol being in error due to a burst of errors. The alarming extension of this is that the next but one FEC symbol being in error has a probability given by:

$$P_{err}(\text{next but one FEC symbol}) = \frac{(0.75^{10} + 0.75^9 + 0.75^8 + 0.75^7 + 0.75^6)}{5} = 0.1086$$

$$P_{err}(\text{next but one FEC symbol}) = 0.75^5 * P_{err}(\text{next but one FEC symbol}) \text{ Note } 0.75^5 = 0.2373$$



# PAM-4 Error propagation options for 1 DFE tap only



0 can only error to 1  
 If correct margin =  $n$   
 If incorrect margin =  $n-1$

1 can error to 0 or 3  
 If correct margin =  $n$   
 If 0 margin = 2  
 If 3 margin =

0 to 0 (prior bit 0)  
 If correct margin = 4  
 If incorrect margin = 6

0 to 0 (prior bit 1)  
 If correct margin = 4  
 If incorrect margin = 6

In 50% of all cases the margin is reduced by  $2 * \text{DFE tap value}$  and could cause an error to propagate